CHOOSING PARAMETERS IN BLOCK-ITERATIVE OR ORDERED SUBSET RECONSTRUCTION

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Viewed abstractly, all the algorithms considered here are designed to provide a nonnegative solution $x$ to the system of linear equations $y = Px$, where $y$ is a vector with positive entries and $P$ a matrix whose entries are nonnegative and with no purely zero columns.

The expectation maximization maximum likelihood (EMML) method in emission tomography [1] and the simultaneous multiplicative algebraic reconstruction technique (SMART) [2, 3, 4, 5] are slow to converge on large data sets; accelerating convergence through the use of block-iterative or ordered subset versions of these algorithms is a topic of considerable interest. These block-iterative versions involve relaxation and normalization parameters the correct selection of which may not be obvious to all users. The algorithms are not faster merely by virtue of being block-iterative; the correct choice of the parameters is crucial. Through a detailed discussion of the theoretical foundations of these methods we come to a better understanding of the precise roles these parameters play.
The notion of cross-entropy or the Kullback-Leibler distance is central to our discussion. For positive numbers $a$ and $b$ let

$$KL(a, b) = a \log(a/b) + b - a;$$

also let $KL(a, 0) = +\infty$ and $KL(0, b) = b$. It is easily seen that $KL(a, b) > 0$ unless $a = b$. We extend this Kullback-Leibler distance component-wise to vectors $x$ and $z$ with nonnegative entries:

$$KL(x, z) = \sum_{j=1}^{J} KL(x_j, z_j).$$

Note that $KL(x, z)$ and $KL(z, x)$ are generally not the same. While the KL distance is not a metric in the usual sense it does have certain properties involving best approximation that are similar to those of the square of the Euclidean metric.

The methods based on cross-entropy, such as the multiplicative version of the algebraic reconstruction technique (ART), the MART [6], its simultaneous version, SMART, the expectation maximization maximum likelihood method (EMML) and all block-iterative versions of these algorithms apply to nonnegative systems that we denote by $Px = y$, where $y$ is a vector of positive entries, $P$ is a matrix with entries $P_{ij} \geq 0$ such that for each $j$ the sum $s_j = \sum_{i=1}^{I} P_{ij}$ is positive and we seek a solution $x$ with nonnegative entries. If no nonnegative $x$ satisfies $y = Px$ we say the system is inconsistent.
Simultaneous iterative algorithms employ all of the equations at each step of the iteration; block-iterative methods do not. For the latter methods we assume that the index set \( \{ i = 1, \ldots, I \} \) is the (not necessarily disjoint) union of the \( N \) sets or blocks \( B_n, n = 1, \ldots, N \). We shall require that \( s_{nj} = \sum_{i \in B_n} P_{ij} > 0 \) for each \( n \) and each \( j \). Block-iterative methods like ART and MART for which each block consists of precisely one element are called row-action or sequential methods. We begin our discussion with the SMART and the EMML method.

Both the SMART and the EMML method provide a solution of \( y = Px \) when such exist and (distinct) approximate solutions in the inconsistent case. Both begin with an arbitrary positive vector \( x^0 \). Having found \( x^k \) the iterative step for the SMART is

**SMART:**

\[
x_j^{k+1} = x_j^k \exp \left( s_j^{-1} \sum_{i=1}^{I} P_{ij} \log \left( \frac{y_i}{(Px^k)_i} \right) \right)
\]  

while that for the EMML method is

**EMML:**

\[
x_j^{k+1} = x_j^k s_j^{-1} \sum_{i=1}^{I} P_{ij} \frac{y_i}{(Px^k)_i}
\]
The following theorems summarize what we know of SMART and EMML.

Theorem 1 \textit{In the consistent case the SMART converges to the unique nonnegative solution of $y = Px$ for which the distance $\sum_{j=1}^{J} s_j KL(x_j, x_j^0)$ is minimized. In the inconsistent case it converges to the unique nonnegative minimizer of the distance $KL(Px, y)$ for which $\sum_{j=1}^{J} s_j KL(x_j, x_j^0)$ is minimized; if $P$ and every matrix derived from $P$ by deleting columns has full rank then there is a unique nonnegative minimizer of $KL(Px, y)$ and at most $I - 1$ of its entries are nonzero.}

Theorem 2 \textit{In the consistent case the EMML algorithm converges to a nonnegative solution of $y = Px$. In the inconsistent case it converges to a nonnegative minimizer of the distance $KL(y, Px)$; if $P$ and every matrix derived from $P$ by deleting columns has full rank then there is a unique nonnegative minimizer of $KL(y, Px)$ and at most $I - 1$ of its entries are nonzero.}

In the consistent case there may be multiple nonnegative solutions and the one obtained using the EMML algorithm will depend on the starting vector $x^0$; how it depends on $x^0$ is an open question. These theorems are special cases of more general results on block-iterative methods.
Those who have used the SMART or the EMML on sizable problems have certainly noticed that they are both slow to converge. An important issue, therefore, is how to accelerate convergence. One popular method is through the use of block-iterative (or ordered subset) methods. To illustrate block-iterative methods and to motivate our subsequent discussion we consider now the ordered subset EM algorithm (OSEM) [7], which is a popular technique in some areas of medical imaging, as well as an analogous version of SMART, which we shall call here the OSSMART. The OSEM algorithm is now used quite frequently in tomographic image reconstruction, where it is acknowledged to produce usable images significantly faster then EMML method.

The idea behind the OSEM (OSSMART) is simple: the iteration looks very much like the EMML (SMART), but at each step of the iteration the summations are taken only over the current block. The blocks are processed cyclically.
The OSEM iteration is the following: for \( k = 0, 1, \ldots \) and \( n \) the index of the current block or subset, having found \( x^k \) let

OSEM:

\[
x^{k+1}_j = x^k_j s_n^{-1} \sum_{i \in B_n} P_{ij} \frac{y_i}{(P x^k)_i}.
\]  
(3)

The OSSMART has the following iterative step:

OSSMART:

\[
x^{k+1}_j = x^k_j \exp(s_n^{-1} \sum_{i \in B_n} P_{ij} \log \frac{y_i}{(P x^k)_i}).
\]  
(4)
In general we do not expect block-iterative algorithms to converge in the inconsistent case, but to exhibit subsequential convergence to a limit cycle. We do, however, want them to converge to a solution in the consistent case; in general, the OSEM and OSSMART do not. These two algorithms are known to converge to a solution in the consistent case when the matrix $P$ and the set of blocks satisfy the condition known as subset balance, which means that the sums $s_{nj}$ depend only on $j$ and not on $n$. While subset balance may be approximately valid in some special cases it is overly restrictive, eliminating, for example, almost every set of blocks whose cardinalities are not all the same. When the OSEM does well in practice in medical imaging it is probably because the $N$ is not large and only a few iterations are carried out.

The experience with the OSEM is encouraging, however, and strongly suggests that an equally fast, but mathematically rigorous, block-iterative version of EMML could be found; this is the *rescaled block-iterative* EMML (RBI-EMML)[8]. Both RBI-EMML and an analogous corrected version of OSSMART, the RBI-SMART, provide fast convergence to a solution in the consistent case, for any choice of blocks.
We consider now block-iterative formulations of the SMART and EMML that are general enough to include all of the variants we wish to discuss. In fact, our initial formulations will be too general and will need to be restricted in certain ways to guarantee and to accelerate convergence.

We begin with the block-iterative version of the SMART, which we shall denote BI-SMART. These methods were known prior to the discovery of RBI-EMML and played an important role in that discovery; the importance of rescaling for acceleration was apparently not appreciated, however. The SMART was discovered in 1972, independently, by Darroch and Ratcliff [2], working in statistics, and by Schmidlin [3] in medical imaging. Block-iterative versions of SMART are also treated in [2], but they also insist on subset balance; the inconsistent case was not considered. We start by considering a formulation of BI-SMART that is general enough to include all of the variants we wish to discuss. As we shall see, this formulation is too general and will need to be restricted in certain ways to obtain convergence.
Initially, we let the BI-SMART iterative step be defined as

\[ x_{j}^{k+1} = x_{j}^{k} \exp \left( \beta_{nj} \sum_{i \in B_{n}} \alpha_{ni} P_{ij} \log \left( \frac{y_{i}}{(P x^{k})_{i}} \right) \right), \quad (5) \]

for \( j = 1, 2, \ldots, J \), \( n = k(\text{mod } N) + 1 \) and \( \beta_{nj} \) and \( \alpha_{ni} \) arbitrary positive weights. Our convergence proof requires that \( \beta_{nj} \) be separable, that is,

\[ \beta_{nj} = \gamma_{j} \delta_{n} \]

for each \( j \) and \( n \) so that (5) becomes

**BI-SMART:**

\[ x_{j}^{k+1} = x_{j}^{k} \exp \left( \gamma_{j} \delta_{n} \sum_{i \in B_{n}} \alpha_{ni} P_{ij} \log \left( \frac{y_{i}}{(P x^{k})_{i}} \right) \right). \quad (6) \]

We also require

\[ \gamma_{j} \delta_{n} \sigma_{nj} \leq 1, \quad (7) \]

for \( \sigma_{nj} = \sum_{i \in B_{n}} \alpha_{ni} P_{ij} \).

With these conditions satisfied we have the following result.

**Theorem 3** Let there be nonnegative solutions of \( y = P x \). For any positive vector \( x^{0} \) and any collection of blocks \( \{B_{n}, n = 1, \ldots, N\} \) the BI-SMART sequence \( \{x^{k}\} \) given by (6) converges to the unique solution of \( y = P x \) for which the weighted cross-entropy, given by \( \sum_{j=1}^{J} \gamma_{j}^{-1} KL(x_{j}, x_{j}^{0}) \), is minimized.
We see from the theorem that how we select the $\gamma_j$ is determined by how we wish to weight the terms in the sum $\sum_{j=1}^{J} \gamma_j^{-1} KL(x_j, x_0^j)$. In some cases we want to minimize the cross-entropy $KL(x, x^0)$ subject to $y = Px$; in this case we would select $\gamma_j = 1$. In other cases we may have some prior knowledge as to the relative sizes of the $x_j$ and wish to emphasize the smaller values more; then we may choose $\gamma_j$ proportional to our prior estimate of the size of $x_j$. Having selected the $\gamma_j$, convergence will be accelerated if we select $\delta_n$ as large as permitted by the condition $\gamma_j \delta_n \sigma_{nj} \leq 1$. This suggests that we take

$$\delta_n = 1/\max\{\sigma_{nj} \gamma_j, j = 1, ..., J\}. \quad (8)$$

The rescaled BI-SMART (RBI-SMART) as presented in [8, 9, 10] uses this choice, but with $\alpha_{ni} = 1$ for each $n$ and $i$.

Let’s look now at some of the other choices for these parameters that have been considered in the literature. First, we notice that the OSSMART does not generally satisfy the requirements, since in (4) the choices are $\alpha_{ni} = 1$ and $\beta_{nj} = s_{nj}^{-1}$; the only times this is acceptable is if the $s_{nj}$ are separable; that is, $s_{nj} = r_j t_n$ for some $r_j$ and $t_n$. This is slightly more general than the condition of subset balance and is sufficient for convergence of OSSMART, since, for $\gamma_j = \alpha_{ni} = 1$ and $\delta_n$ as in (8), the BI-SMART reduces to the OSSMART.
In [4] Censor and Segman make the choices $\beta_{nj} = 1$ and $\alpha_{ni} > 0$ such that $\sigma_{nj} \leq 1$ for all $n$ and $j$. In those cases in which $\sigma_{nj}$ is much less than 1 for each $n$ and $j$ their iterative scheme is probably excessively relaxed; it is hard to see how one might improve the rate of convergence by altering only the weights $\alpha_{ni}$, however. Limiting the choice to $\gamma_j \delta_n = 1$ reduces our ability to accelerate this algorithm.

The original SMART in (1) uses $N = 1$, $\gamma_j = s_j^{-1}$ and $\alpha_{ni} = \alpha_i = 1$. Clearly (7) is satisfied; in fact it becomes an equality now. For the row-action version of SMART, the multiplicative ART (MART), due to Gordon, Bender and Herman [6], we take $N = I$ and $B_n = B_i = \{i\}$ for $i = 1, ..., I$.

Darroch and Ratcliff included a discussion of a block-iterative version of SMART in their 1972 paper [2]. Close inspection of their version reveals that they require that $s_{nj} = \sum_{i \in B_n} P_{ij} = 1$ for all $j$. Since this is unlikely to be the case initially, we might try to rescale the equations or unknowns to obtain this condition. However, unless $s_{nj} = \sum_{i \in B_n} P_{ij}$ depends only on $j$ and not on $n$, which is the subset balance property used in [7], we cannot redefine the unknowns in a way that is independent of $n$. 

12
The MART begins with a strictly positive vector $x^0$ and has the iterative step

**MART:**

$$x_{j}^{k+1} = x_{j}^{k} \left( \frac{y_i}{(P \cdot x^k)_i} \right)^{m_i - 1} P_{ij},$$

(9)

for $j = 1, 2, ..., J$, $i = k \text{(mod I)} + 1$ and $m_i > 0$ chosen so that $m_i^{-1} P_{ij} \leq 1$ for all $j$. Convergence of the MART is generally faster for smaller $m_i$, so a good choice is $m_i = \max \{P_{ij}, j = 1, ..., J\}$. Although this particular choice for $m_i$ is not explicitly mentioned in the various discussions of MART, it was used in implementations of MART from the beginning.

The MART fails to converge in the inconsistent case. What is always observed, but for which no proof exists, is that, for each fixed $i = 1, 2, ..., I$, as $m \to +\infty$, the MART subsequences $\{x^{mI+i}\}$ converge to separate limit vectors, say $x^{\infty,i}$. This limit cycle $LC = \{x^{\infty,i} | i = 1, ..., I\}$ reduces to a single vector whenever there is a nonnegative solution of $y = Px$. The greater the minimum value of $KL(Px, y)$ the more distinct from one another the vectors of the limit cycle are. An analogous result is observed for BISMART.
We turn now to the block-iterative version of the EMML algorithm.

Initially, we let the iterative step of the BI-EMML be defined as

\[
x^{k+1}_j = x^k_j (1 - \beta_{nj} \sigma_{nj}) + x^k_j \beta_{nj} \sum_{i \in B_n} \alpha_{ni} P_{ij} \frac{y_i}{(P x^k)_i},
\]

for \( j = 1, 2, ..., J \), \( n = k (\text{mod} \ N) + 1 \) and \( \beta_{nj} \) and \( \alpha_{ni} \) positive. As in the case of BI-SMART, our convergence proof requires that \( \beta_{nj} \) be separable, that is,

\[
\beta_{nj} = \gamma_j \delta_n
\]

for each \( j \) and \( n \) and that

\[
\gamma_j \delta_n \sigma_{nj} \leq 1,
\]

for \( \sigma_{nj} = \sum_{i \in B_n} \alpha_{ni} P_{ij} \). The BI-EMML then becomes

BI-EMML:

\[
x^{k+1}_j = x^k_j (1 - \gamma_j \delta_n \sigma_{nj}) + x^k_j \gamma_j \delta_n \sum_{i \in B_n} \alpha_{ni} P_{ij} \frac{y_i}{(P x^k)_i},
\]

With these conditions satisfied we have the following result.

Theorem 4 Let there be nonnegative solutions of \( y = Px \). For any positive vector \( x^0 \) and any collection of blocks \( \{ B_n, n = 1, ..., N \} \) the BI-EMML sequence \( \{ x^k \} \) given by (11) converges to a nonnegative solution of \( y = Px \).
When there are multiple nonnegative solutions of $y = Px$ the solution obtained by BI-EMML will depend on the starting point $x^0$, but precisely how it depends on $x^0$ is an open question. Also, in contrast to the case of BI-SMART, the solution can depend on the particular choice of the blocks.

Having selected the $\gamma_j$, convergence will be accelerated if we select $\delta_n$ as large as permitted by the condition $\gamma_j \delta_n \sigma_{nj} \leq 1$. This suggests that once again we take $\delta_n$ as in (8). The rescaled BI-EMML (RBI-EMML) as presented in [8, 9, 10] uses this choice, but with $\alpha_{ni} = 1$ for each $n$ and $i$.

Let’s look now at some of the other choices for these parameters that have been considered in the literature. First, we notice that the OSEM does not generally satisfy the requirements, since in (3) the choices are $\alpha_{ni} = 1$ and $\beta_{nj} = s_{nj}^{-1}$; the only times this is acceptable is if the $s_{nj}$ are separable; that is, $s_{nj} = r_j t_n$ for some $r_j$ and $t_n$. This is slightly more general than the condition of subset balance and is sufficient for convergence of OSEM, since, for $\gamma_j = \alpha_{ni} = 1$ and $\delta_n$ as in (8), the BI-EMML reduces to the OSEM.
The original EMML in (2) uses $N = 1$, $\gamma_j = s_j^{-1}$ and $\alpha_{ni} = \alpha_i = 1$. Clearly (7) is satisfied; in fact it becomes an equality now. Notice that the calculations required to perform the BI-SMART are somewhat more complicated than those needed in BI-EMML. Because the MART converges rapidly in most cases there is considerable interest in the row-action version of EMML. It was clear from the outset that using the OSEM in a row-action mode does not work. We see from the formula for BI-EMML that the proper row-action version of EMML, which we call the EM-MART, has the iterative step

**EM-MART:**

$$x_j^{k+1} = (1 - \delta_i \gamma_j \alpha_{ii} P_{ij}) x_j^k + \delta_i \gamma_j \alpha_{ii} P_{ij} \frac{y_i}{(P x^k)_i},$$  \hspace{1cm} (12)

with

$$\gamma_j \delta_i \alpha_{ii} P_{ij} \leq 1$$

for all $i$ and $j$.  

16
The optimal choice would seem to be to take $\delta_i \alpha_{ii}$ as large as possible; that is, to select

$$\delta_i \alpha_{ii} = 1/\max \{\gamma_j P_{ij}, j = 1, ..., J\}.$$  

With this choice the EM-MART is called the rescaled EM-MART (REM-MART). The EM-MART fails to converge in the inconsistent case. What is always observed, but for which no proof exists, is that, for each fixed $i = 1, 2, ..., I$, as $m \to +\infty$, the EM-MART subsequences $\{x_{mI+i}\}$ converge to separate limit vectors, say $x^{\infty,i}$. This limit cycle $\text{LC} = \{x^{\infty,i} | i = 1, ..., I\}$ reduces to a single vector whenever there is a non-negative solution of $y = Px$. The greater the minimum value of $KL(y, Px)$ the more distinct from one another the vectors of the limit cycle are. An analogous result is observed for BI-EMML.
We must mention a method that closely resembles the REM-MART, the row-action maximum likelihood algorithm (RAMLA), which was discovered independently by Browne and De Pierro [11]. The RAMLA avoids the limit cycle in the inconsistent case by using strong underrelaxation involving a decreasing sequence of relaxation parameters $\lambda_k$. The RAMLA has the following iterative step:

RAMLA:

$$x_j^{k+1} = (1 - \lambda_k \sum_{i \in B_n} P_{ij})x_j^k + \lambda_k x_j^k \sum_{i \in B_n} P_{ij} \left( \frac{y_i}{(P x^k)_i} \right), \quad (13)$$

where the positive relaxation parameters $\lambda_k$ are chosen to converge to zero and $\sum_{k=0}^{+\infty} \lambda_k = +\infty$.

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The results discussed in this poster are taken from a paper to be published in the IEEE Transactions on Image Processing. The full paper is available on my website (http://faculty.uml.edu/cbyrne/cbyrne.html).
References


