Toward a Reliable Data Transport Architecture for Optical Burst-Switched Networks

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Presentation Outline

- Introduction to Optical Transport Paradigms
- Optical Burst Switching
- · Reliable Data Transport in OBS
 - Loss Minimization Mechanisms
 - Loss Recovery Mechanisms
- · Conclusion and Future Work



Applications Demands

Service Optical Transport Applications Requirements **Paradigms Voice Over IP Streaming Video High Bandwidth Optical Circuit Switching Grid Computing Dynamic Provisioning Optical Packet Switching Storage Area Networks** Reliability **Optical Burst Switching Multimedia Low Latency Data**

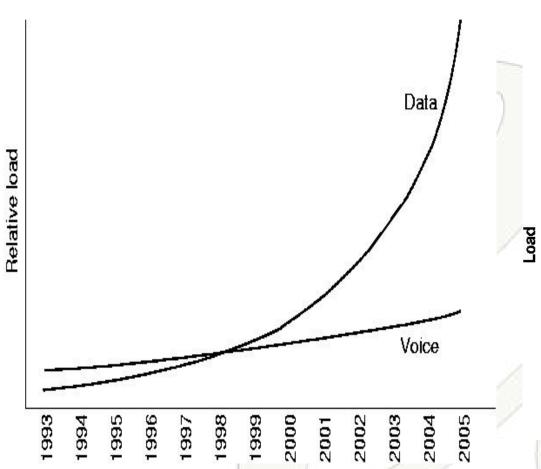


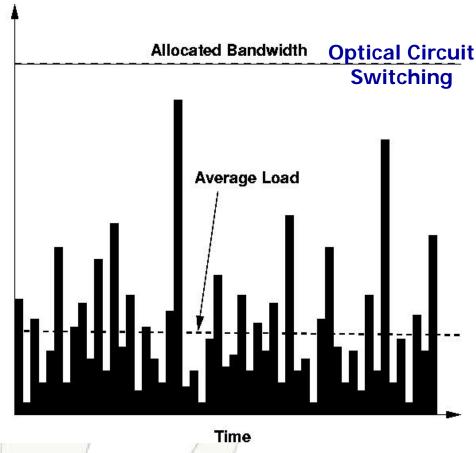
Optical Circuit Switching

- · For each request,
 - Set-up a static circuit (lightpath)
 - Transfer data
 - Release connection
- Pros:
 - Suitable for smooth, longer-term, high-bandwidth applications
- · Cons:
 - Long circuit set-up latency
 - Inefficient for short-term bursty applications



Optical Circuit Switching (cont.)





- Circuit switched networks optimized for Voice
- Data: Accounts for majority of total traffic
- Data tends to be bursty
- Static bandwidth allocation is **not** efficient



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Optical Packet Switching

- A photonic packet contains header and payload
- · Header is processed all-optically at each node

· Pros:

- Statistical multiplexing of data
- Suitable for bursty traffic

Cons:

- Very fast switching speeds (nanoseconds)
- Synchronization



Optical Burst Switching

- Multiple IP packets assembled into a burst
- An out-of-band control header transmitted ahead of each data burst

Pros:

- Statistical multiplexing of data
- Suitable for bursty traffic
- Low data-transfer latency
- Electronic control plane (practically feasible)
- Optical data plane (high-speed)



Motivation for OBS

Optical Switching Paradigm	Bandwidth Utilization	Setup Latency	Switching Speed Req.	Proc. / Sync. Overhead	Traffic Adaptively
Optical Circuit Switching	Low	High	Slow	Low	Low
Optical Packet Switching	High	Low	Fast	High	High
Optical Burst Switching	High	Low	Medium	Low	High

OBS combines the best of the two while avoiding their shortcomings

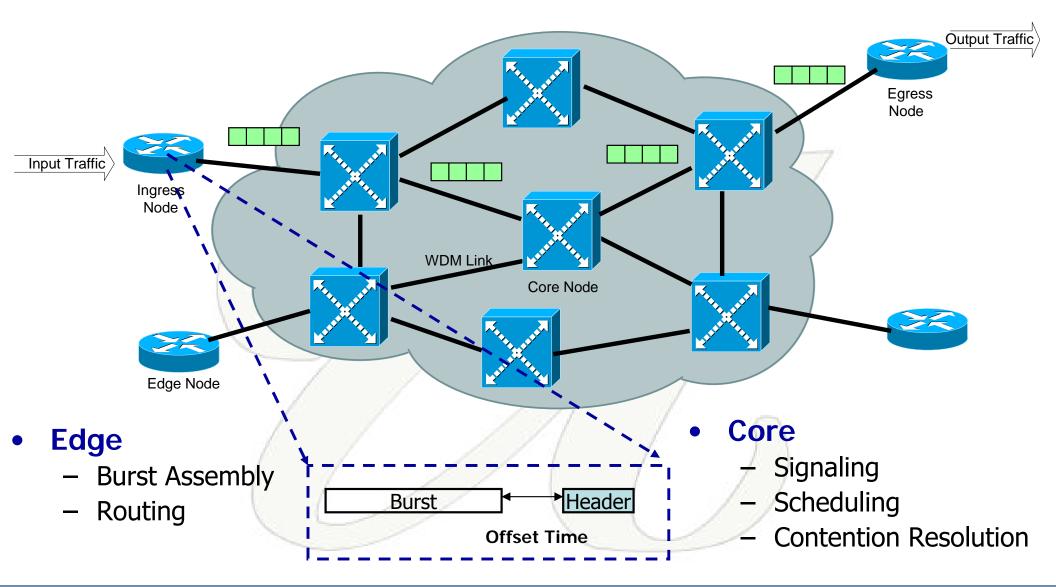


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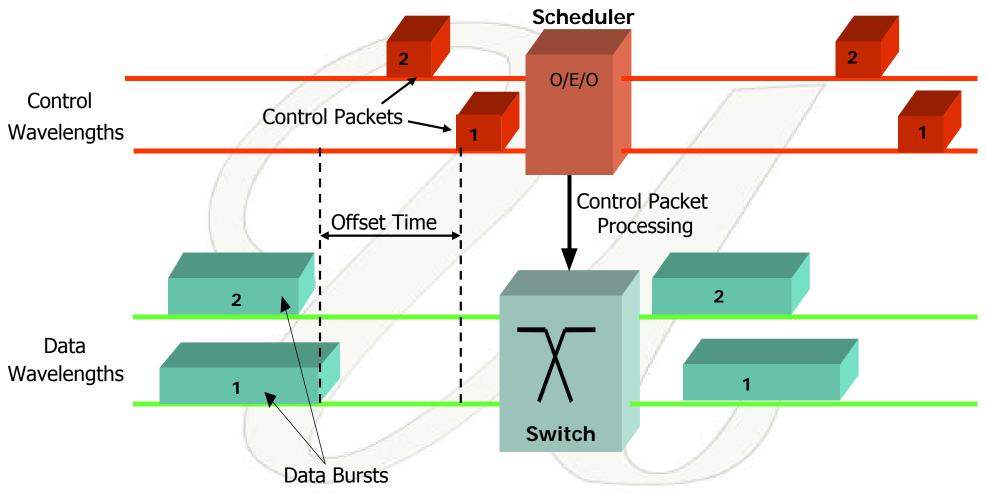


OBS Network Architecture





OBS Node Architecture



Adopted from Qiao



Burst Assembly

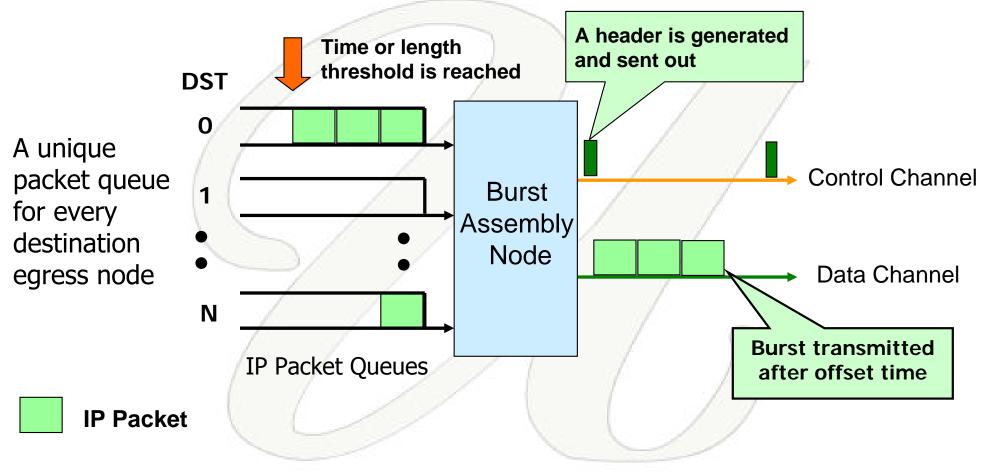
- Aggregate multiple (IP) packets going to the same destination into a single burst
- · Assembly Mechanisms: Timer-based and Threshold-based

- Timer-based assembly:
 - After a fixed timer interval, all the packets in the queue are framed into a single burst
- Threshold-based assembly:
 - After a fixed length threshold is reached, all the packets in the queue are framed into a single burst.



Burst Assembly

Aggregate multiple (IP) packets going to the same destination into a single burst

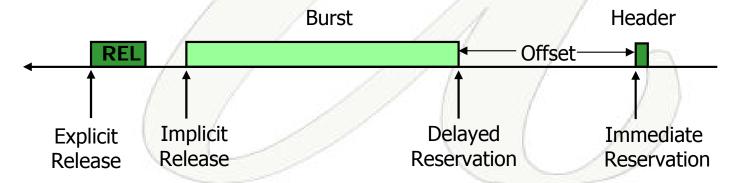


Adopted from Qiao



Signaling Technique

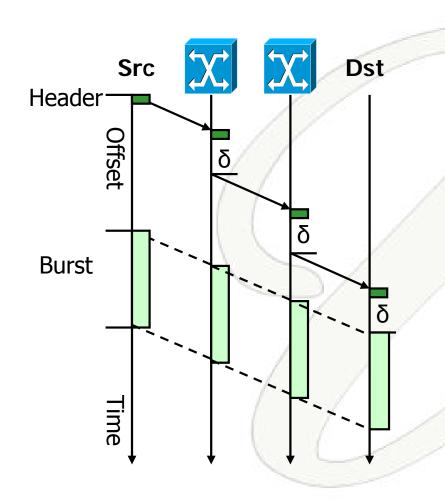
- One-way based (un-acknowledged) signaling
- Reservation Mechanism: Based on the start of the reservation
 - Immediate Reservation: Immediately after the control heater
 - Delayed Reservation: At the start of the burst
- Release Mechanism: Based on the release of the reservation
 - Implicit Release: based on burst length information
 - Explicit Release: explicit release control packet used



Tradeoff: Efficiency vs. Simplicity



Just-Enough-Time (JET) Signaling



- Delayed Reservation and Implicit Release
- Header contains burst length, offset time, source, destination
- Offset time necessary for processing of header at intermediate nodes without buffering the data burst
- Just-In-Time (JIT): Immediate Reservation and Explicit Release

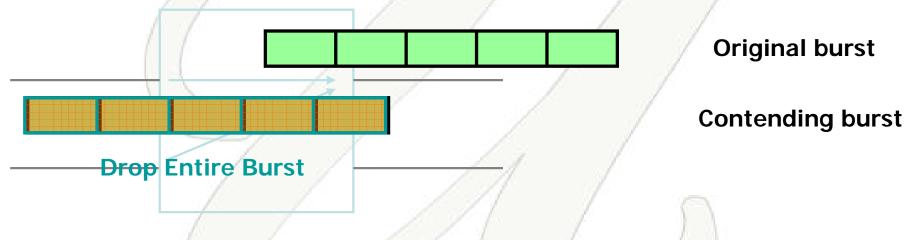


Data Loss in OBS: Burst Contentions

•Contention occurs when more than one burst attempts to go out of same output port (or wavelength) at the same time

·Unique to all-optical networks

- -Traditional networks employ electronic buffering to resolve contentions
- -Lack of optical buffers (cannot store light)



Core Switch

- Drop Policy:
 - One of the bursts will be dropped in its entirety
 - Even though overlap between the bursts may be minimal



TCP over OBS

- Transmission Control Protocol (TCP)
 - Majority of applications depend on TCP for reliable data tfr.
 - TCP assumes packet loss is always due to network congestion
 - TCP congestion avoidance mechanisms reduce sending rate in the event of a packet loss

· OBS

- Random burst loss occurs even when network is NOT congested
- TCP over OBS
 - False Timeout time out when network is NOT congested
 - TCP falsely reduces send rate even when network is NOT congested
 - Significantly degrades throughput of high-bandwidth apps

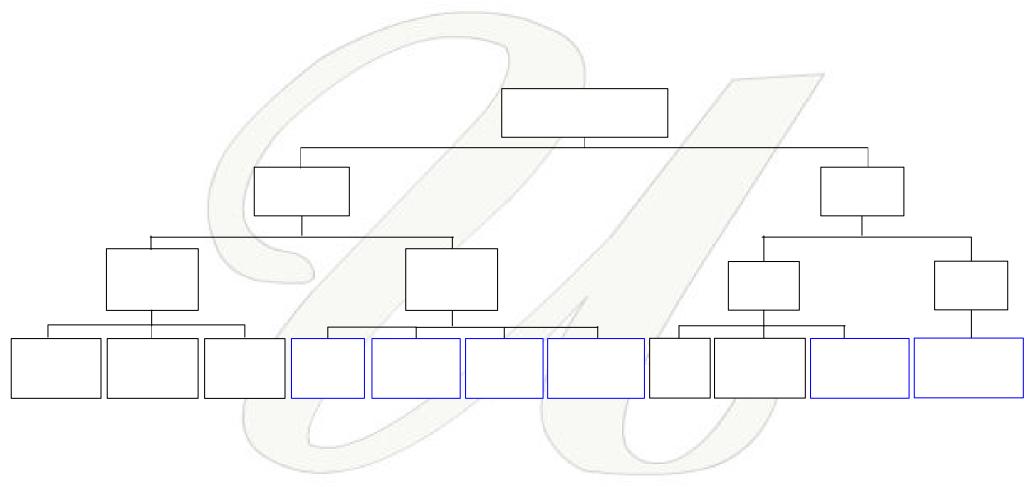


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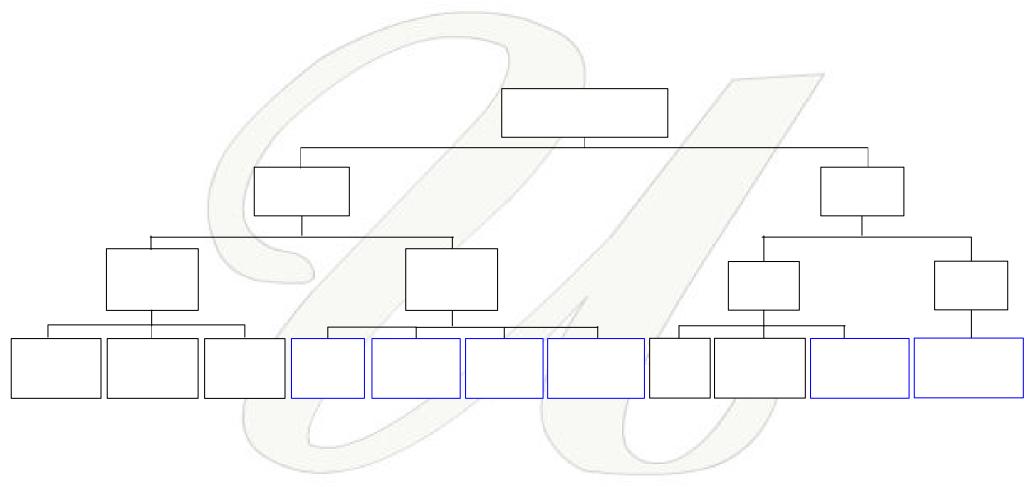
Toward a Reliable OBS







Toward a Reliable OBS







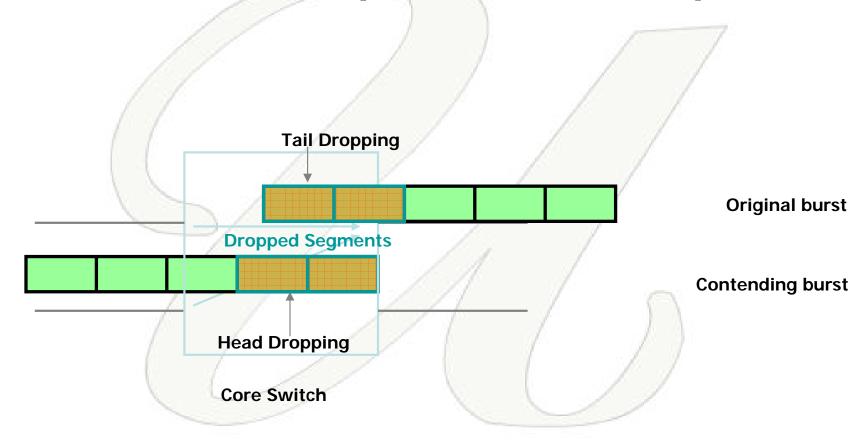
Traditional Contention Resolution

- Optical Buffering (FDLs)
 - Achieved through Fiber Delay Lines
 - Issues: Limited buffer capacity and additional hardware cost
- Wavelength Conversion
 - Converting the wavelength of an incoming channel to another wavelength at the outgoing channel
 - Issues: Additional hardware cost
- Deflection Routing
 - Deflect contending bursts to alternate port
 - No additional hardware cost
 - Issues: Higher delay and out-of-sequence delivery



Burst Segmentation

- When contention occurs, only overlapping segments are dropped
- Two Approaches: Head Dropping and Tail Dropping
- Details: Vokkarane and Jue [IEEE ICC 2002, New York]





Evaluation Criteria

- Evaluation of proposed policies
 - Average end-to-end packet loss probability
 - Average end-to-end packet delay (hops)
 - TCP throughput
- Numerical Analysis
 - Analytical modeling Markov models
 - Simulation results Discrete-event simulations



Burst Segmentation: Analytical Loss Model

- Burst Arrivals: Poisson Process
- M/G/1/1 Queueing Model

Burst Length Distribution (After k hops):
$$G_{l_k^{sd}}(t) = 1 - (1 - G_{l_{k-1}^{sd}}(t))(e^{-\lambda_{l_k^{sd}}})$$

$$= 1 - (1 - G_{l_0^{sd}}(t))e^{-\left(\sum_{i=1}^k \lambda_{l_i^{sd}}\right)t}.$$

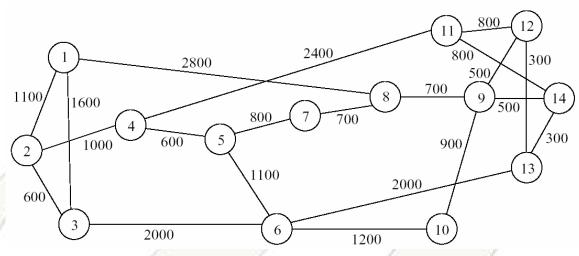
End-to-End Packet Loss:

$$P_{\text{loss}} = \sum_{s} \sum_{d} \frac{\lambda^{sd}}{\lambda} P_{\text{loss}}^{sd}.$$

[Vokkarane: IEEE JSAC 2003, SPIE Optical Networks 2003]



Simulation Network



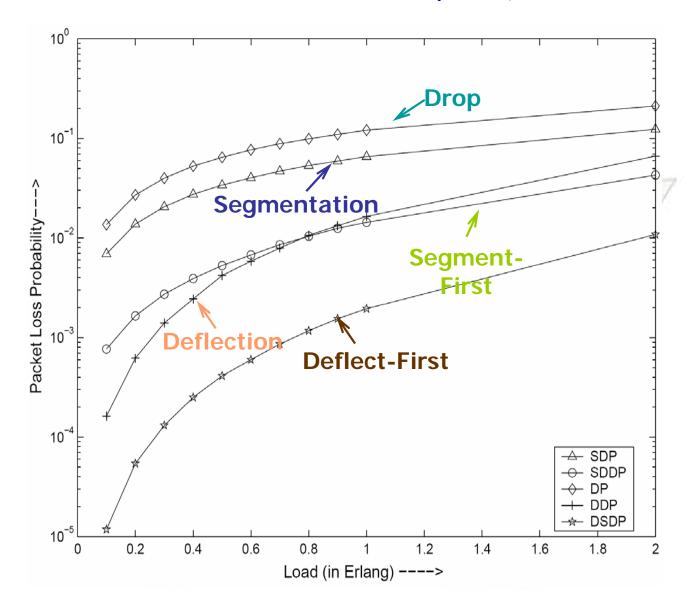
Assumptions

14-node NSFNET

Burst Arrivals	Poisson		
Average Burst Length	100 µs (exponentially dist.)		
Link Transmission Rate	10 Gb/s		
Packet Length	1500 Bytes		
Switching Time	10 µs		
Optical Buffering	NO		
Wavelength Conversion	NO		

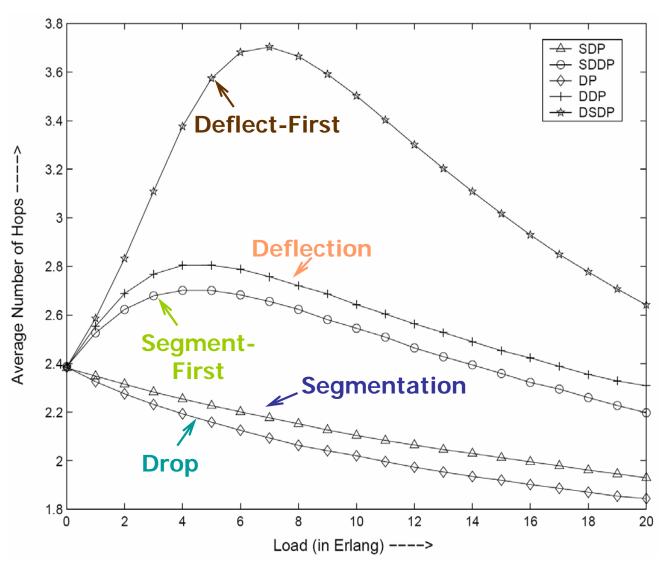


Packet Loss Performance



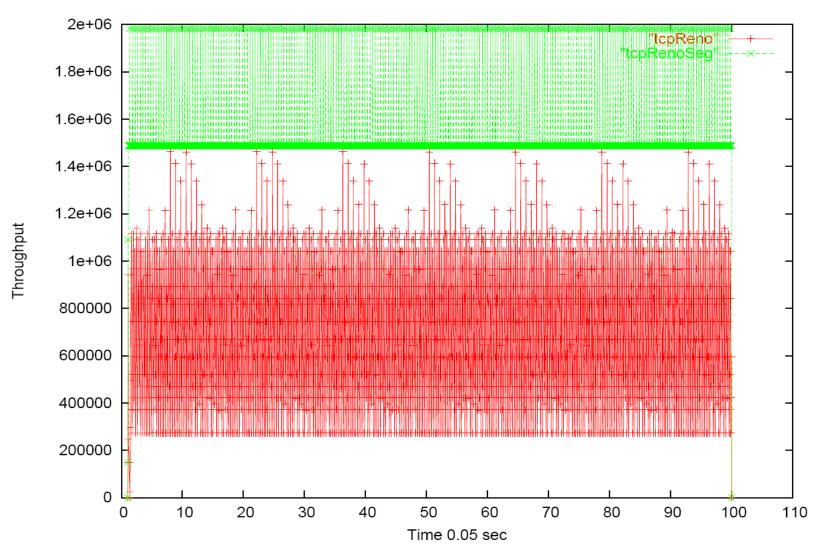


Average Number of Hops (~Delay)



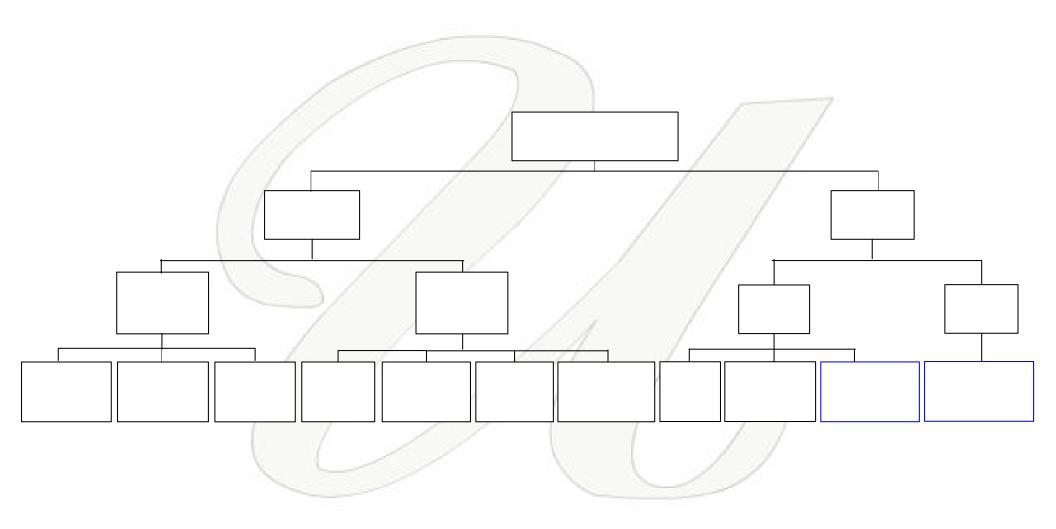


TCP Throughput





Toward a Reliable OBS





Reactive Loss Recovery: Burst Retransmission

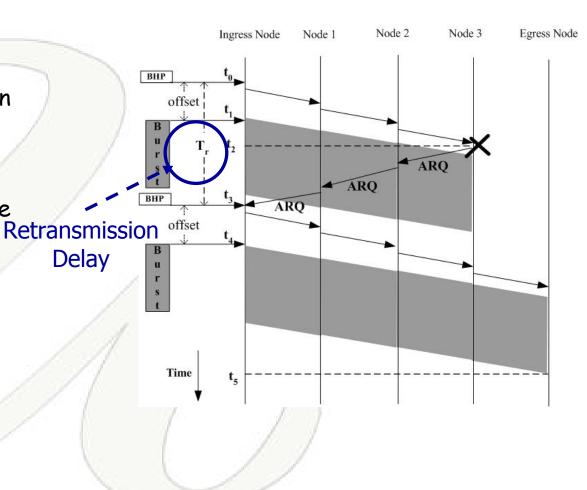
Objective

 To recover from burst loss when network is not congested

Basic idea

Retransmit lost bursts at source nodes

Stop retransmission when
 Tr > Delay Constraint (δ)



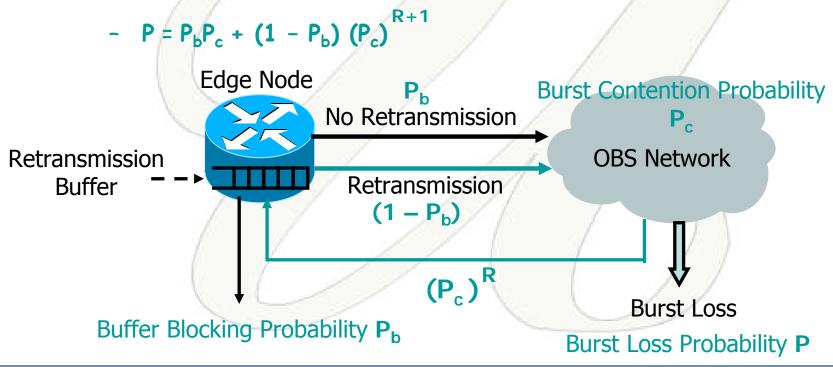
[Zhang and Vokkarane: IEEE GLOBECOM 2005]



Analysis for Burst Retransmission

Objective

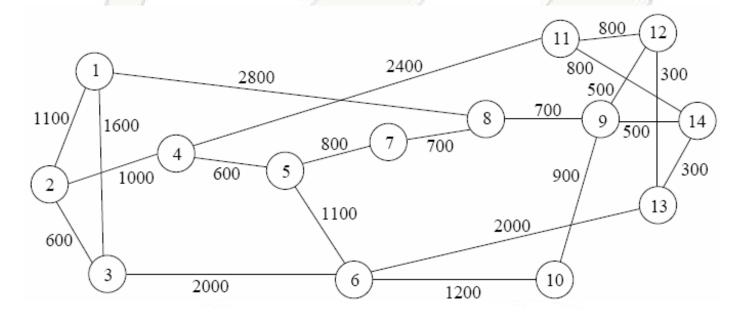
- Analyze the average burst loss probability in the network Basic idea
 - No retransmission for bursts blocked by retransmission buffers





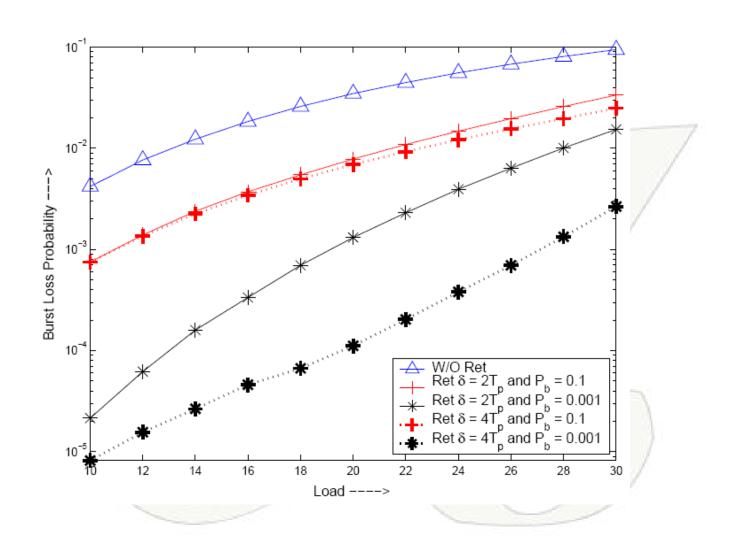
Simulation Assumptions

- No. of wavelengths on each link is 4
- Transmission rate on a wavelength is 10 Gb/s
- Burst arrival is Poisson
- · Traffic are uniformly distributed
- Average burst length are 100 µs



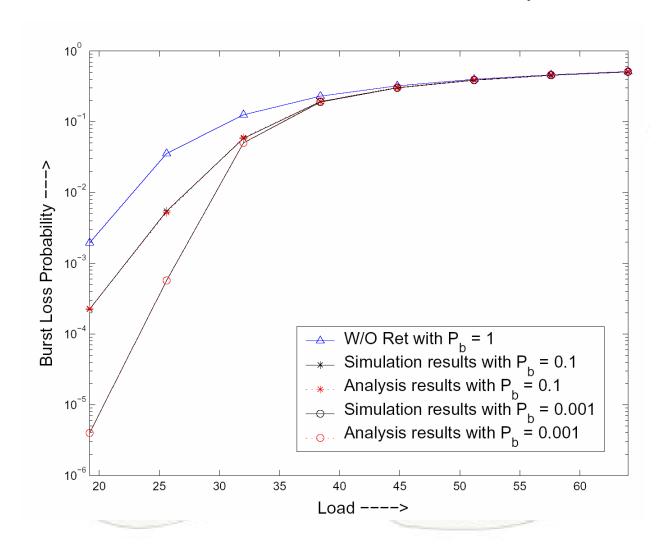


Burst Loss Probability



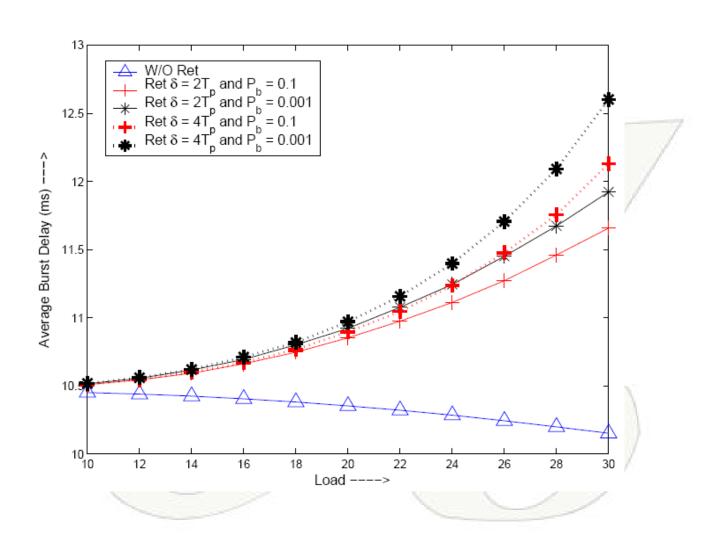


Analysis and Simulation Results for Burst Loss Probability



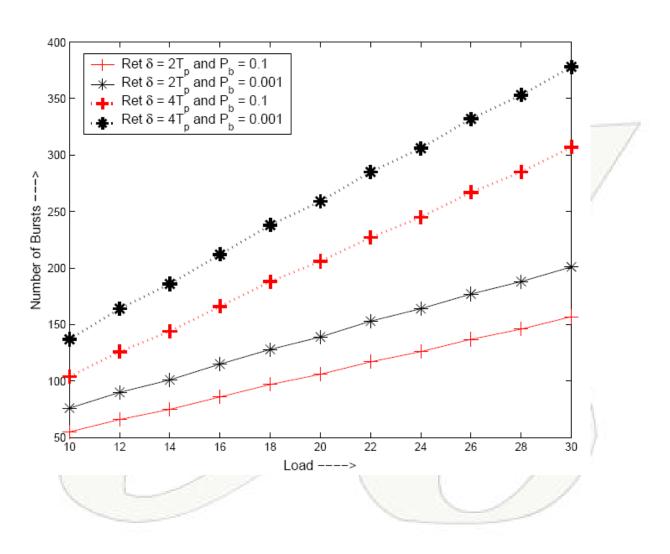


Average Burst Delay



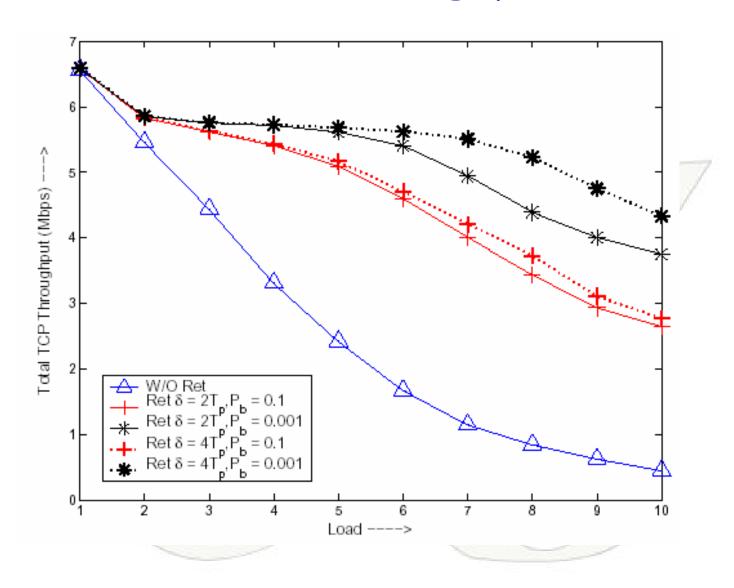


Buffer Capacity at Edge Nodes



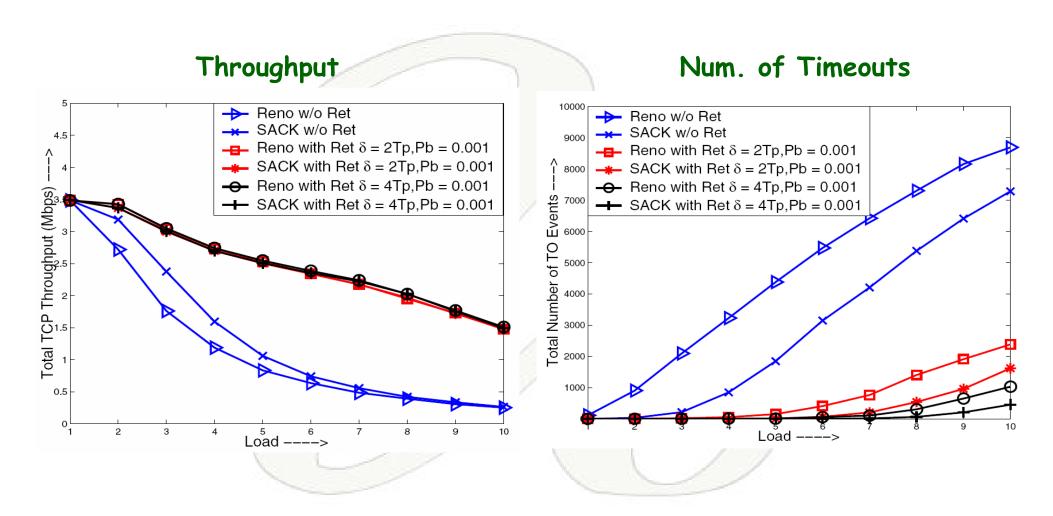


TCP Throughput





Performance of TCP Versions





Burst Retransmission

Pros

- Reduce burst loss probability
- Correctly indicate network congestion (min FTOs)
- Significantly improve TCP throughput

· Cons

- Additional electronic buffers at edge nodes
- Longer delay for retransmitted bursts



Proactive Loss Recovery: Forward Redundancy (FR)

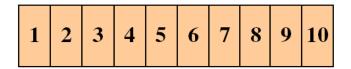
- Some or all the original packets of a burst are copied and sent in the forward direction from source to destination
- Receiver can recover from selective packet loss in the forward direction

FR Policies

- Serial or Parallel FR
- Partial (< 100%) or Complete (>= 100%) FR
- Same or Disjoint path Protection



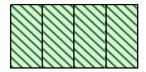
Serial Forward Redundancy (SFR)





i) Serial Forward Redundancy





ii) Parallel Forward Redundancy



Loss-Sensitive Packet



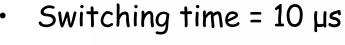
Redundant Packet

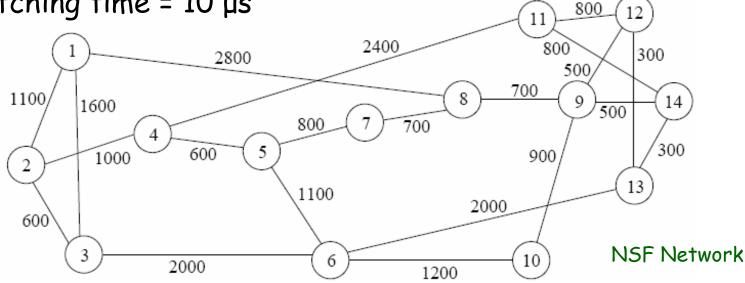
We have evaluated SFR [IEEE GridNets 2005, IEEE WOCN 2006]



Simulation Assumptions

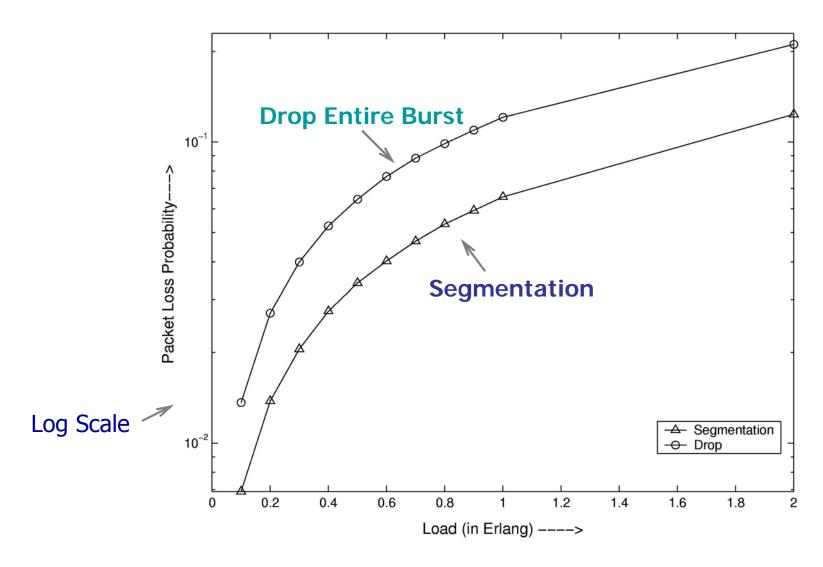
- No. of wavelengths on each link: 8
- Transmission rate: 10 Gbps per wavelength
- Burst arrivals: Poisson
- Packet size: 1250 byte [10 Kb]
- Fixed burst length: 100 packets [1Mb]
- Traffic: uniformly distributed





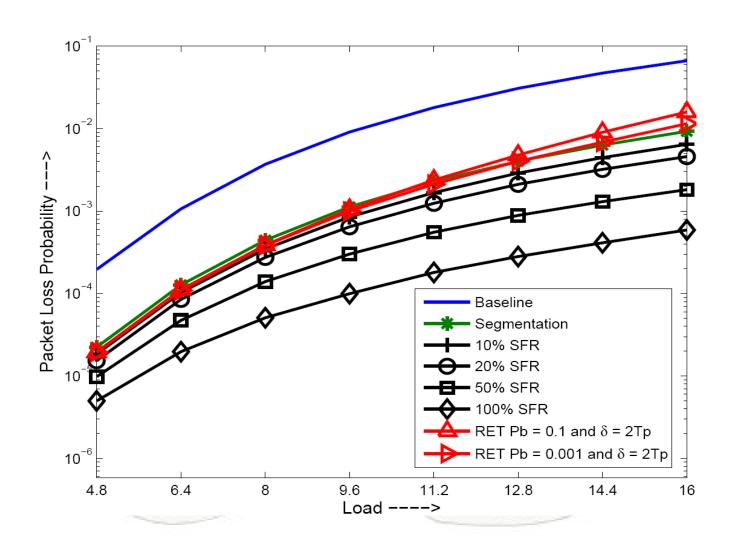


Packet Loss Performance



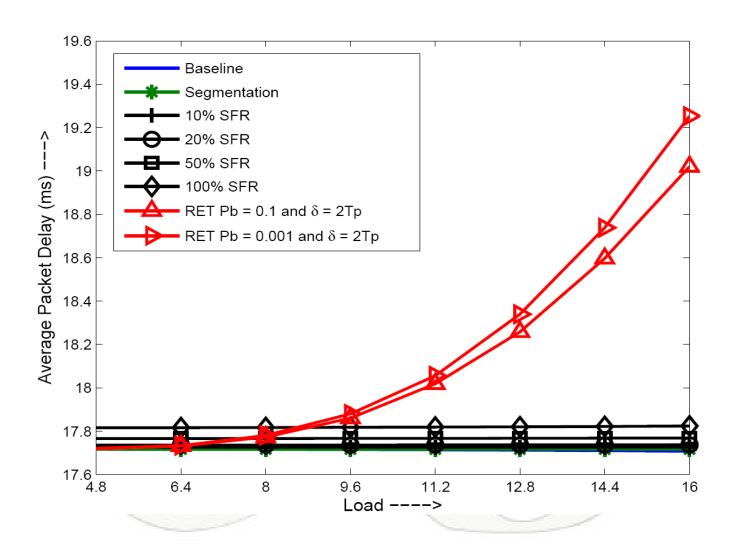


Packet Loss Probability



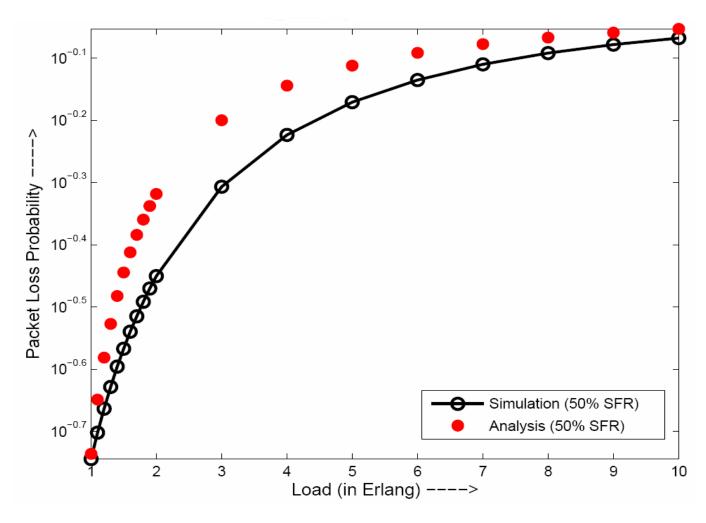


End-to-End Packet Delay



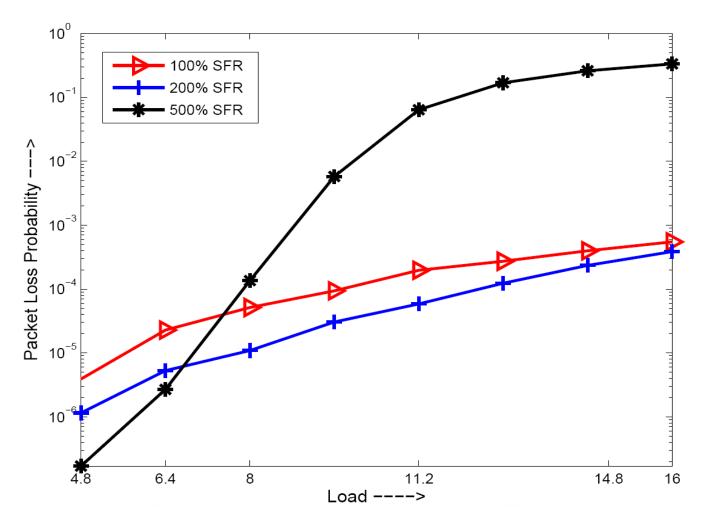


Analytical Loss Model Results



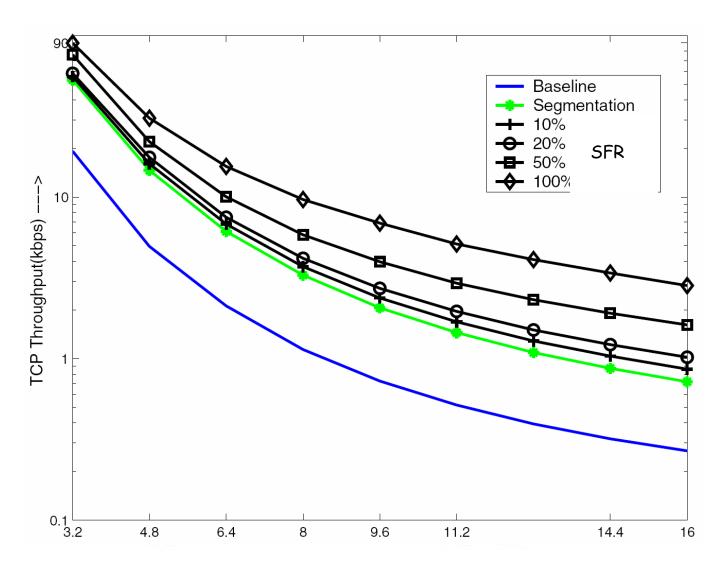


Packet Loss Probability





TCP Throughput





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Conclusion

OBS Network

- Promising optical core data-transport paradigm
- Suited for delay-sensitive applications

Loss Minimization and Loss Recovery Mechanisms

- Evaluated several new mechanisms
- Proposed mechanisms significantly improves the reliability of data transfer over OBS networks

Future Work

- Develop dynamic mechanisms
- Impact on newer high-speed TCP versions



Thank You

http://www.cis.umassd.edu/~vvokkarane/

